

### Three-Valued Logics

Notation: p, q, r atomic formulas; A, B, C arbitrary formulas

Values = {T, F, N}

Valuation = assignment of a value from Values to each atomic formula

Can be extended to an assignment of values to each formula by inductive definition using truth tables...

Truth Tables:

Bochvar = Weak Kleene = Halldén

$\sim$		$\&$	T	N	F	$\vee$	T	N	F	$\supset$	T	N	F	$\equiv$	T	N	F
T	F	T	T	N	F	T	T	N	T	T	T	N	F	T	T	N	F
N	N	N	N	N	N	N	N	N	N	N	N	N	N	N	N	N	N
F	T	F	F	N	F	F	T	N	F	F	T	N	T	F	F	N	T

Strong Kleene = Körner

$\sim$		$\&$	T	N	F	$\vee$	T	N	F	$\supset$	T	N	F	$\equiv$	T	N	F
T	F	T	T	N	F	T	T	T	T	T	T	N	F	T	T	N	F
N	N	N	N	N	F	N	T	N	N	N	T	N	N	N	N	N	N
F	T	F	F	F	F	F	T	N	F	F	T	T	T	F	F	N	T

Łukasiewicz

$\sim$		$\&$	T	N	F	$\vee$	T	N	F	$\supset$	T	N	F	$\equiv$	T	N	F
T	F	T	T	N	F	T	T	T	T	T	T	N	F	T	T	N	F
N	N	N	N	N	F	N	T	N	N	N	T	<b>T</b>	N	N	N	<b>T</b>	N
F	T	F	F	F	F	F	T	N	F	F	T	T	T	F	F	N	T

'Determinately true' operator

$d$	
T	T
N	F
F	F

Three ways to define validity:

1. preservation of truth (Łukasiewicz, Bochvar, Kleene, Tye):  
 $\Gamma \models B$  iff for every valuation  $v$ , if every formula in  $\Gamma$  is T on  $v$ , then B is T on  $v$   
 corresponds to designating T
2. preservation of non-falsity (Halldén):  
 $\Gamma \models B$  iff for every valuation  $v$ , if every formula in  $\Gamma$  is T or N on  $v$ , then B is T or N on  $v$   
 corresponds to designating both T and N
3. preservation of truth and non-falsity:  
 $\Gamma \models B$  iff for every valuation  $v$ , if every formula in  $\Gamma$  is T on  $v$ , then B is T on  $v$ , *and*  
 if every formula in  $\Gamma$  is T or N on  $v$ , then B is T or N on  $v$   
 corresponds to preservation of degree (in ordering  $T > N > F$ )

## Continuum-valued Logics

Values =  $[0,1]$  (set of reals from 0 to 1 inclusive)

Valuation = an assignment of values from Values to atomic formulas.

Semantics for connectives:

Łukasiewicz

$$|\sim A| = 1 - |A|$$

$$|A \vee B| = \max(|A|, |B|)$$

$$|A \& B| = \min(|A|, |B|)$$

$$|A \supset B| = \begin{cases} 1 & \text{if } |A| \leq |B| \\ 1 - (|A| - |B|) & \text{if } |A| > |B| \end{cases}$$

$$|A \equiv B| = 1 - (\max(|A|, |B|) - \min(|A|, |B|))$$

Goguen

$$|A \& B| = |A| \times |B|$$

$$|A \supset B| = \begin{cases} 1 & \text{if } |A| \leq |B| \\ |B| / |A| & \text{if } |A| > |B| \end{cases}$$

Edgington

non-degree functional proposal

$$|\sim A| = 1 - |A|$$

$$|A \& B| = |A| \times |B| \text{ given } A$$

$$|A \vee B| = |A| + |B| - |A \& B|$$

Three ways to define validity:

1. preservation of perfect truth (Williamson, Edgington):

$\Gamma \models B$  iff for every valuation  $v$ , if every formula in  $\Gamma$  is 1 on  $v$ , then  $B$  is 1 on  $v$   
corresponds to designating 1

2. preservation of degree of truth (Machina):

$\Gamma \models B$  iff for every valuation  $v$ , the value of  $B$  on  $v \geq$  the lowest value assigned to a member of  $\Gamma$  on  $v$   
corresponds to ordering the values and saying that a valid argument can't take you lower

3. conclusion can have no more falsity than the premises combined:

$\Gamma \models B$  iff for every valuation  $v$ ,  $1 - |B|_v \leq \sum_{A \in \Gamma} (1 - |A|_v)$

cf. Edgington p. 307, but note that this is not how she *defines* validity. (She uses (1).)